

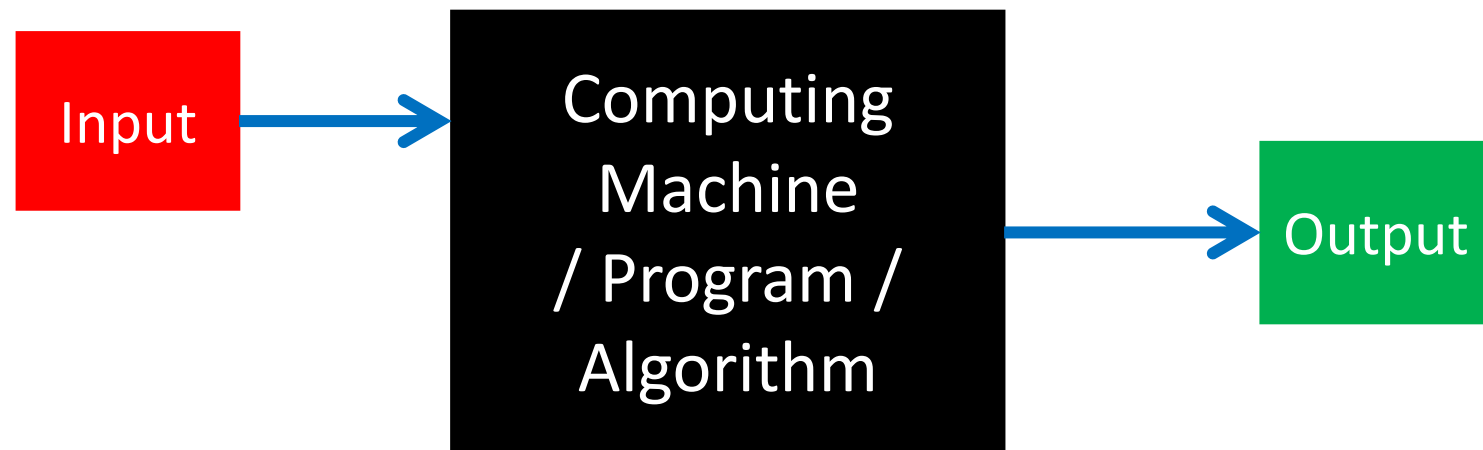
CS3102

September 6

Questions

Recall: “Carnot Engine” of computing

- What “type” is the input?
 - String (bit strings)
- What “type” is the output?
 - String
- What have we put in the black box so far?
 - Circuits of logic gates
 - Python/Java/C++
 - X8086
 - NAND-CIRC



What goes in here?

What do we compute on?

- Always strings! (that represent other things)
 - Images
 - DNA
 - Web pages
 - Temperature
- To compute on non-string things, we need a *representation scheme*
 - Function $E: U \rightarrow \{0,1\}^*$
 - E should be one-to-one

Model of Computing

- To define a model of computing, we need:
 - A way to give it input
 - Pre-defining the values of the “input” gates
 - A way to get output out of it
 - Pre-defined some gates as “output” gates, and when execution stopped, their values defined the output
 - How do we go from input to output?
 - Execution model: deterministic (non-random)
 - For any gate where all of its incoming wires had defined values, its value becomes the result of applying a particular operation to the values its incoming wires
 - Representation
 - Wires, gates, gates have labels (AON: AND/OR/NOT/INPUT/OUTPUT; NAND:NAND/INPUT/OUTPUT)

Showing AON=NAND

- Anything I can do with AON I can also do with NAND
 - Vice-versa
- Two models of computing are “equivalent” if every function you can implement with one of them you can also implement with the other
- Find a way of converting instances of one model into instances of the other such that we compute the same function
- To go from AON to NAND:
 - Do AND with NANDs
 - Do OR with NANDs
 - NOT with NANDs
- To go from NAND to AON:
 - Do NAND with ANDs/Ors/NOTs

Syntactic Sugar

- What is it?
 - Adding in capability of naming subroutines to make your model more human-readable
- Why use it?
 - Readability
 - Similar to relationship between assembly named memory locations vs. explicit
 - Succinctness b/c you can reuse “code”

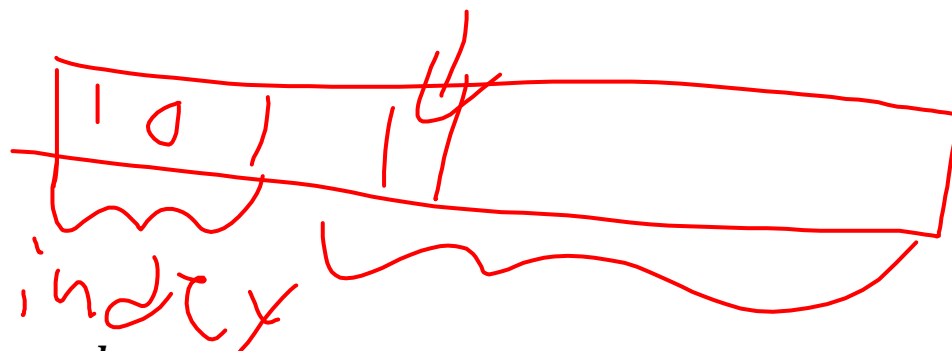
Lookup_k and *Add_n*

- What does the subscript mean?

- What the size of the input is

- k is the number of bits in the index (total of $2^k + k$)

- n is the number of bits in each number we're adding together (total of $2n$)



$$\text{Compare}_k: \{0,1\}^{2k} \rightarrow \{0,1\}$$

- Using only NAND gates, write a function to compare two-bit integers
 - $\text{Compare}_2(a_1, a_0, b_1, b_0) = 1$ if integer $a_1a_0 >$ integer b_1b_0
 - $\text{Compare}_2(1100) = 1$
 - $\text{Compare}_2(0101) = 0$
- $\text{Compare}_1(a, b) =$

a	b	$c, (a, b)$
0	0	0
0	1	0
1	0	1
1	1	0

$$\text{Compare}_1(a, b) = \neg(a \rightarrow b) \equiv a \wedge \neg b$$

*Compare*₁(a, b):

notb = NAND(b, b)

anb = NAND(a, notb)

notanb = NAND(anb, anb)

$Compare_k(s_a, s_b)$:

$eq0 = \text{NOT}(\text{XOR}(s_a[0], s_b[0]))$ # 3 gates

$compkm1 = Compare_{k-1}(s_a[1:], s_b[1:])$ # C_{k-1}

$comp1 = Compare_1(s_a[0], s_b[0])$ # 3

return IF(eq0, compkm1, comp1) # 4 gates

If a more significant bit differs between the two, then we know the answer, otherwise we need to check more bits

$$C_k = C_{k-1} + 10$$
$$C_1 = 3$$

- $k = 1$
 - 3 gates
- $k = 2$
 - 13 gates
- $k = 3$
 - 23 gates
- $k = 4$
 - 33 gates
- ...
- In general for k
 - $3 + 10(k - 1)$

Induction to show: $C_k \leq 3 + 10(k - 1)$

- Base case: $k = 1$
 - $C_1 \leq 3 + 10(1 - 1)$
 - $3 \leq 3$
- Inductive step: to show if $C_k \leq 3 + 10(k - 1)$ then $C_{k+1} \leq 3 + 10((k + 1) - 1)$
 $C_{k+1} = C_k + 10$

$$C_k \leq 3 + 10(k - 1)$$

$$C_{k+1} \leq 3 + 10(k - 1) + 10$$

$$C_{k+1} \leq 3 + 10k$$

$$C_{k+1} \leq 3 + 10((k + 1) - 1)$$